

# A Note on Carnap's Continuum and the Weak State Description Analogy Principle

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## Abstract

Under the assumption of Atom Exchangeability and Regularity we give a characterization of the probability functions in Carnap's Continuum of Inductive Methods in terms of the Weak State Description Analogy Principle.

Key words: State Description Analogy Principle, Carnap's Continuum, Pure Inductive Logic, Analogy, Rationality.

## The Characterization

We shall assume the context and notation of Unary Pure Inductive Logic as set out in [2]. Throughout  $w$  will be a probability function on the set  $SL$  of sentences of the unary language  $L$  with  $q \geq 2$  predicate symbols which satisfies Ex.

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For atoms

$$\alpha_j = \bigwedge_{i=1}^q R_i^{\epsilon_i}(x), \quad \alpha_k = \bigwedge_{i=1}^q R_i^{\delta_i}(x)$$

set

$$\Delta(\alpha_r, \alpha_m) = \{i \mid \epsilon_i = \delta_i\}.$$

### The Weak State Description Analogy Principle, WSDAP

For atoms  $\alpha_r, \alpha_m, \alpha_k$ , if  $\Delta(\alpha_r, \alpha_m) \subseteq \Delta(\alpha_r, \alpha_k)$  then

$$w(\alpha_r(a_{n+2}) \mid \alpha_k(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{h_i}(a_i)) \leq w(\alpha_r(a_{n+2}) \mid \alpha_m(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{h_i}(a_i)). \quad (1)$$

**Theorem 1.** *Let  $w$  satisfy Atom Exchangeability and Regularity. Then  $w$  satisfies WSDAP just if  $w = c_\lambda^L$  for some  $0 < \lambda \leq \infty$ .*

*Proof.* First suppose that  $w$  satisfies WSDAP. Let  $r, m, k$  be distinct and such that

$$\Delta(\alpha_r, \alpha_m) \subseteq \Delta(\alpha_r, \alpha_k).$$

By WSDAP we must have

$$w(\alpha_r(a_{n+2}) \mid \alpha_m(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{h_i}(a_i)) \leq w(\alpha_r(a_{n+2}) \mid \alpha_k(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{h_i}(a_i)). \quad (2)$$

But we can (for  $q \geq 2$ ) produce a permutation  $\sigma$  of  $\{1, 2, \dots, 2^q\}$  which leaves  $r$  fixed and is such that

$$\Delta(\alpha_r, \alpha_{\sigma(m)}) \supset \Delta(\alpha_r, \alpha_{\sigma(k)})$$

whilst by using Atom Exchangeability, Ax, on (2) we also have that

$$w(\alpha_r(a_{n+2}) \mid \alpha_{\sigma(m)}(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{\sigma(h_i)}(a_i)) \leq w(\alpha_r(a_{n+2}) \mid \alpha_{\sigma(k)}(a_{n+1}) \wedge \bigwedge_{i=1}^n \alpha_{\sigma(h_i)}(a_i)).$$

This is only consistent with WSDAP if we actually have equality in (2).

Indeed by the transitivity of equality we must also have equality in (2) for *any*  $\alpha_m, \alpha_k$  not equal to  $\alpha_r$ . From this it follows that for state descriptions

$$\Theta = \bigwedge_{i=n+1}^m \alpha_{h_i}(a_i), \quad \Phi = \bigwedge_{i=n+1}^m \alpha_{g_i}(a_i)$$

where  $r$  does not appear amongst the  $h_i$  or  $g_i$ , we must have that

$$w(\alpha_r(a_1) \mid \bigwedge_{i=2}^n \alpha_r(a_i) \wedge \Theta) = w(\alpha_r(a_1) \mid \bigwedge_{i=2}^n \alpha_r(a_i) \wedge \Phi). \quad (3)$$

For if this failed then, by changing one atom at a time and using Ex, it would have to fail in a case when just  $h_{n+1} \neq g_{n+1}$  and  $h_j = g_j$  for  $n+2 \leq j \leq m$ , which would contradict equality holding in (2). But, under Ax, (3) is just Johnson's Sufficientness Postulate whose only solutions  $w$  for  $q \geq 2$  and  $w$  regular are Carnap's  $c_\lambda^L$  for  $0 < \lambda \leq \infty$ .

In the other direction it is straightforward to check by direct calculation that these  $c_\lambda^L$  satisfy WSDAP.  $\square$

Unfortunately, as pointed out in [3], in the case of the  $c_\lambda^L$  ( $\lambda > 0$ ) WSDAP as above holds with equality in (2) so it can hardly be said to capture more than a faint hint of analogy. From other work in this area, see for example [1], it is clear that this collapse is very much a consequence of assuming Ax, much more structure remains if instead we only assume Px+SN.

## References

- [1] D'Asaro, F.A., *Analogical Reasoning in Unary Inductive Logic*, Master's Thesis, University of Manchester, 2014. Available at <http://www.maths.manchester.ac.uk/~jeff/theses/fdathesis.pdf>
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