## MATH43001/63001, January 2011 Exam, Solutions<sup>1</sup>

- **A1.** (i)  $f(x_1, f(x_1, x_2)) \in TL$  since  $x_1, x_2 \in TL$  by Te1, so  $f(x_1, x_2) \in TL$  by Te2 and  $f(x_1, f(x_1, x_2)) \in TL$  by Te2 again.
- (ii)  $f((f(x_1, x_2), x_1) \notin TL)$  since this word has different numbers of right and left round brackets and we can prove by induction on |t| that any  $t \in TL$  has the same number. [Not necessary to give the proof but for the record: Clearly true if t is a constant or free variable  $x_i$  (when there are zero of either) and if  $t = f(t_1, \ldots, t_n)$  then the number of '(' in t equals 1 plus the number in  $t_1, \ldots, t_n$ , equals 1 plus the number of ')' in  $t_1, \ldots, t_n$ , by inductive hypothesis, equals the number of ')' in t.]
- (iii)  $\forall w_1 \neg R(w_1, x_1) \in FL$  since  $R(x_2, x_1) \in FL$  by L1, so  $\neg R(x_2, x_1) \in FL$  by L2, and finally then  $\forall w_1 \neg R(w_1, x_1) \in FL$  by L3.
- (iv)  $\forall w_1 \neg R(w_2, x_1) \notin FL$  since we can prove by induction on  $|\theta|$  for  $\theta \in FL$  that if  $w_2$  occurs in  $\theta$  then so does either  $\exists w_2$  or  $\forall w_2$ , which rules out  $\forall w_1 \neg R(w_2, x_1)$  being in FL. [Again no need to prove this but for the record: Clearly true, vacuously, for  $R(t_1, t_2)$ , and if it holds for  $\phi$ ,  $\psi$  then it holds for  $\neg \phi$ ,  $(\phi \land \psi)$ ,  $(\phi \lor \psi)$ ,  $(\phi \to \psi)$ . Also if it holds for  $\eta$  and  $\eta$  does not mention  $w_j$ , if  $j \neq 2$  then it holds for  $\exists w_j \eta(w_j/x_i)$  and  $\forall w_j \eta(w_j/x_i)$ , whilst if j = 2 then the condition holds trivially for  $\exists w_2 \eta(w_2/x_i)$  and  $\forall w_2 \eta(w_2/x_i)$ .]

(v) 
$$M \models \forall w_1 \forall w_2 (R(w_1, w_2) \rightarrow R(w_2, w_1)) \iff$$
 for all  $n, m \in \mathbb{N}^+$ , if  $n < m$  then  $m < n$ ,

which is false since, e.g. 1 < 2 but  $2 \nleq 1$ .

(vi) 
$$M \models \exists w_1 \forall w_2 \neg R(w_2, f(w_1, w_2)) \iff$$
  
there is an  $n \in \mathbb{N}^+$  such that for all  $m \in \mathbb{N}^+$ ,  $m \not< nm$ ,

which is true when we take n = 1 since  $m \not< 1 \times m$  for any  $m \in \mathbb{N}^+$ .

(vii) 
$$M \models \forall w_1 (R(w_1, f(w_1, w_1)) \rightarrow \forall w_2 R(w_2, f(w_1, w_2))) \iff$$
 for all  $n \in \mathbb{N}^+$ , if  $n < n^2$  then for all  $m \in \mathbb{N}^+$ ,  $m < nm$ .

This is true since if  $n \in \mathbb{N}^+$  and  $n < n^2$  then n > 1 so m < nm for  $m \in \mathbb{N}^+$ .

$$\theta_{1}(x_{1}, x_{2}) = R(f(x_{1}, x_{1}), x_{2})$$

$$\theta_{2}(x_{1}, x_{2}) = (\neg R(x_{1}, x_{2}) \land \neg R(x_{2}, x_{1}))$$

$$\theta_{3}(x_{1}, x_{2}) = (R(x_{1}, x_{2}) \land \neg \exists w_{1} (R(x_{1}, w_{1}) \land R(w_{1}, x_{2})))$$

$$\theta_{4}(x_{1}, x_{2}) = \exists w_{1} \theta_{2}(f(x_{1}, w_{1}), w_{2}) = \exists w_{1} (\neg R(f(x_{1}, w_{1}), x_{2}) \land \neg R(x_{2}, f(x_{1}, w_{1})))$$

$$\phi = \forall w_{1} \exists w_{2} R(w_{1}, f(w_{1}, w_{2})) \text{ (since this fails in } K \text{ when } w_{1} = 0).$$

<sup>&</sup>lt;sup>1</sup>These solutions are more detailed than I would expect in the exam. That's because I want them to also serve an educational purpose when given with 'last year's paper' next year(!)

**A2.** A suitable logical equivalent (there are many possibilities here) in PNF is

$$\forall w_2 \, \forall w_1 \, (P(w_2) \to \neg R(w_1)).$$

It is enough to just write this down for the marks but for the record we could argue:

$$\neg \exists w_1 R(w_1) \equiv \forall w_1 \neg R(w_1) \text{ and } \exists w_1 P(w_1) \equiv \exists w_2 P(w_2)$$

by the 'Useful Equivalents' (UEs for short).

$$\therefore (\exists w_1 P(w_1) \to \neg \exists w_1 R(w_1)) \equiv (\exists w_2 P(w_2) \to \forall w_1 \neg R(w_1)) \text{ by Lemma 1,}$$

$$\therefore (\exists w_1 P(w_1) \to \neg \exists w_1 R(w_1)) \equiv \forall w_2 (P(w_2) \to \forall w_1 \neg R(w_1))$$

by UEs and transitivity of  $\equiv$ . Also by UEs,

$$(P(x_2) \rightarrow \forall w_1 \neg R(w_1)) \equiv \forall w_1 (P(x_2) \rightarrow \neg R(w_1))$$

so by Lemma 1,

$$\forall w_2 (P(w_2) \rightarrow \forall w_1 \neg R(w_1)) \equiv \forall w_2 \forall w_1 (P(w_2) \rightarrow \neg R(w_1))$$

and the result follows by transitivity of  $\equiv$ .

**A3.** A formal proof of  $\exists w_1 \, \theta(w_1) \to \phi \vdash \forall w_1 \, (\theta(w_1) \to \phi)$  where  $w_1$  does not occur in  $\phi$ :

- 1  $\theta(x_1), \exists w_1 \theta(w_1) \to \phi \mid \exists w_1 \theta(w_1) \to \phi$  REF
- $2 \quad \theta(x_1), \ \exists w_1 \theta(w_1) \to \phi \mid \theta(x_1)$  REF
- $\exists \theta(x_1), \exists w_1 \theta(w_1) \to \phi \mid \exists w_1 \theta(w_1)$   $\exists I, 2$
- $4 \quad \theta(x_1), \exists w_1 \theta(w_1) \rightarrow \phi \mid \phi$  MP, 1, 3
- 5  $\exists w_1 \theta(w_1) \to \phi \mid (\theta(x_1) \to \phi)$  IMR, 4
- 6  $\exists w_1 \theta(w_1) \to \phi \mid \forall w_1 (\theta(w_1) \to \phi) \quad \forall I, 5$

**A4.** Completeness Theorem: For  $\Gamma \subseteq FL$  and  $\theta \in FL$ ,  $\Gamma \vdash \theta \iff \Gamma \models \theta$ . (a) Let M be the structure for L such that  $|M| = \mathbb{N}$ ,  $P^M = \{ n \in \mathbb{N} \mid n \text{ is even } \}$ ,  $Q^M = \{ n \in \mathbb{N} \mid n \text{ is odd } \}$ . Then  $M \models \forall w_1 P(w_1) \rightarrow \forall w_1 Q(w_1) \text{ since } M \nvDash \forall w_1 P(w_1)$ . However  $M \nvDash \forall w_1 (P(w_1) \rightarrow Q(w_1)) \text{ since } 0 \in \mathbb{N} \text{ is even but not odd. Hence}$ 

$$\forall w_1 P(w_1) \rightarrow \forall w_1 Q(w_1) \not\vDash \forall w_1 (P(w_1) \rightarrow Q(w_1))$$

and by the Completeness Theorem

$$\forall w_1 P(w_1) \rightarrow \forall w_1 Q(w_1) \not\vdash \forall w_1 (P(w_1) \rightarrow Q(w_1)).$$

(b) Let M be a structure for L and suppose that

$$M \models \forall w_1 \forall w_2 \left( P(w_1) \lor Q(w_2) \right) \quad \star$$

but

$$M \nvDash \forall w_1 P(w_1) \vee \exists w_2 Q(w_2)$$
 †

Then

$$M \nvDash \forall w_1 P(w_1)$$
 and  $M \nvDash \exists w_2 Q(w_2)$ .

Hence for some  $a \in |M|$ ,  $M \nvDash P(a)$  and also  $M \nvDash Q(a)$  since  $M \nvDash \exists w_2 Q(w_2)$ . Hence  $M \nvDash P(a) \lor Q(a)$ . But this contradicts  $\star$ . Hence given  $\star \dagger$  must fail, so

$$\forall w_1 \forall w_2 \left( P(w_1) \vee Q(w_2) \right) \models \forall w_1 P(w_1) \vee \exists w_2 Q(w_2)$$

and by the Completeness Theorem

$$\forall w_1 \forall w_2 \left( P(w_1) \vee Q(w_2) \right) \vdash \forall w_1 P(w_1) \vee \exists w_2 Q(w_2).$$

- **A5.** (i)+(ii)  $\nvDash$  (iii): Let M be the structure for L such that  $|M| = \mathbb{N}$  and  $R^M = \{\langle n, m \rangle \in \mathbb{N}^2 \mid n < m \}$ . Then (i) is true in M since for every  $n \in \mathbb{N}$  there is  $m \in \mathbb{N}$  such that n < m and (ii) is true in M since  $0 \in \mathbb{N}$  and  $m \not< 0$  for every  $m \in \mathbb{N}$ . However  $M \models R(0,1)$  since 0 < 1 but there is no  $n \in \mathbb{N}$  such that  $M \models R(0,n) \land R(n,1)$ , i.e. 0 < n < 1 so (iii) fails in M.
- (i)+(iii)  $\nvDash$  (ii): Let M be the structure for L with  $|M| = \mathbb{R}$  and  $R^M = \{\langle n, m \rangle \in \mathbb{R}^2 \mid n < m \}$ . Then (i) is true in M since for every  $r \in \mathbb{R}$  there is an  $s \in \mathbb{R}$  such that r < s and (iii) is true in M since if  $r, s \in \mathbb{R}$  and r < s then there is a  $t \in \mathbb{R}$  (for example (r+s)/2) such that r < t < s. However (ii) fails in M since otherwise there would have to be some  $r \in \mathbb{R}$  such that for all  $s \in \mathbb{R}$ ,  $s \not< r$ , which is false (take s = r 1).
- (ii)+(iii)  $\nvDash$  (i): Let M be the structure for L with  $|M| = \{0\}$  and  $R^M = \emptyset$ . Then for any  $s \in |M|$ ,  $\langle 0, s \rangle \notin R^M$  so (ii) holds in M. Also since  $\langle s, r \rangle \notin R^M$  for any  $r, s \in |M|$ ,  $M \nvDash R(s, r)$  and  $M \models R(s, r) \to \exists w_3 (R(w_1, w_3) \land R(w_3, w_2))$ . Hence (iii) holds in M. However (i) fails in M since for the only element of |M|, 0, there is no  $s \in |M|$  such that  $M \models R(0, s)$ , i.e.  $\langle 0, s \rangle \in R^M$ .
- **B6.** Claim For any  $\phi(\vec{x}) \in FL$  and any  $\vec{a} \in |M|$ ,

$$M^* \models \phi(\vec{a}) \iff M \models \phi^*(\vec{a})$$

where (as expected)  $\phi^*(\vec{x})$  is the result of replacing the relation symbol P everywhere in  $\phi(\vec{x})$  by Q.

The claim is proved by induction on  $|\phi|$  (for all  $\vec{a}$  simultaneously). If  $\phi(\vec{x}) = R(x_{i_1}, \ldots, x_{i_m})$  and  $R \neq P$  then  $\phi^*(\vec{x}) = \phi(\vec{x})$  and

$$M \models \phi^*(\vec{a}) \iff M \models \phi(\vec{a}) \iff \langle a_{i_1}, \dots, a_{i_m} \rangle \in R^M$$

$$\iff \langle a_{i_1}, \dots, a_{i_m} \rangle \in R^{M^*} \iff M^* \models \phi^*(\vec{a}) \iff M^* \models \phi(\vec{a}).$$

If R = P then

$$M \models \phi^*(\vec{a}) \iff M \models Q(a_{i_1}, \dots, a_{i_m}) \iff \langle a_{i_1}, \dots, a_{i_m} \rangle \in Q^M$$
  
$$\iff \langle a_{i_1}, \dots, a_{i_m} \rangle \in P^{M^*} \iff M^* \models P(a_{i_1}, \dots, a_{i_m}) \iff M^* \models \phi(\vec{a}).$$

Assuming the result for  $\psi(\vec{x}), \eta(\vec{x}), \chi(x_i, \vec{x})$  (and noticing that  $((\psi(\vec{x}) \wedge \eta(\vec{x})))^* = (\psi^*(\vec{x}) \wedge \eta^*(\vec{x}))$  we have

$$M \models \psi^*(\vec{a}) \land \eta^*(\vec{a}) \iff M \models \psi^*(\vec{a}) \text{ and } M \models \eta^*(\vec{a}) \iff$$

 $\iff M^* \models \psi(\vec{a}) \text{ and } M^* \models \eta(\vec{a}) \text{ (by Ind.Hyp.)} \iff M^* \models \psi(\vec{a}) \land \eta(\vec{a})$ and similarly for the other connectives. Also (noticing that  $(\exists w_j \chi(w_j, \vec{x}))^* = \exists w_j \chi^*(w_j, \vec{x})$ )

$$M \models \exists w_i \, \chi^*(w_i, \vec{a}) \iff \text{ for some } b \in |M|, M \models \chi^*(b, \vec{a}) \iff$$

 $\iff$  for some  $b \in |M| (= |M^*|)$ ,  $M^* \models \chi(b, \vec{a})$  (by Ind.Hyp.  $\iff M^* \models \exists w_j \chi(w_j, \vec{a})$ , and similarly for  $\chi(x_i, \vec{x})$ , completing the induction.

Now suppose that  $\models \theta(\vec{x})$ . Then for any structure M for L and assignment  $\vec{x} \mapsto \vec{a}$ ,  $M^* \models \theta(\vec{a})$  so by the claim  $M \models \theta^*(\vec{a})$ . Hence  $\models \theta^*(\vec{x})$ , as required.

The converse is not true, for example  $\models Q(x_1) \lor \neg Q(x_1)$  but  $\not\models Q(x_1) \lor \neg P(x_1)$ .

**B7.** A proof of EqL(=),  $\forall w_1 R(w_1, w_1) \vdash x_1 = x_2 \to R(x_1, x_2)$ .

1 
$$x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ x_1 = x_2$$
 REF  
2  $| x_1 = x_1$  Eq1  
3  $x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ (x_1 = x_1 \land x_1 = x_2)$  AND, 1, 2  
4  $| \forall w_1, w_2, w_3, w_4 \ ((w_1 = w_3 \land w_2 = w_4) \rightarrow (R(w_1, w_2) \rightarrow R(w_3, w_4)))$  Eq4  
5  $| \forall w_2, w_3, w_4 \ ((x_1 = w_3 \land w_2 = w_4) \rightarrow (R(x_1, w_2) \rightarrow R(w_3, w_4)))$   $\forall O, A$   
6  $| \forall w_3, w_4 \ ((x_1 = w_3 \land x_1 = w_4) \rightarrow (R(x_1, w_2) \rightarrow R(x_1, w_4)))$   $\forall O, B$   
7  $| \forall w_4 \ ((x_1 = x_1 \land x_1 = w_4) \rightarrow (R(x_1, x_1) \rightarrow R(x_1, w_4)))$   $\forall O, B$   
8  $| ((x_1 = x_1 \land x_1 = x_2) \rightarrow (R(x_1, x_1) \rightarrow R(x_1, x_2)))$   $\forall O, B$   
9  $| x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ (R(x_1, x_1) \rightarrow R(x_1, x_2))$  MP, 3, 8  
10  $| x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ R(x_1, x_1)$  REF  
11  $| x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ R(x_1, x_1)$   $\forall O, B$   
12  $| x_1 = x_2, \ \forall w_1 R(w_1, w_1) \ | \ R(x_1, x_2)$  MP, 9, 11  
13  $| \forall w_1 R(w_1, w_1) \ | \ x_1 = x_2 \rightarrow R(x_1, x_2)$  IMR, 12

**B8.** The Compactness Theorem: For L a language and  $\Gamma \subseteq FL$ ,  $\Gamma$  is satisfiable iff every finite subset of  $\Gamma$  is satisfiable.

Suppose on the contrary that there was such a sentence  $\theta$ . Let  $\Gamma$  be the set of sentences  $\{\theta\} \cup \{\neg \phi_n \mid n \in \mathbb{N}^+\}$  of L where  $\phi_n$  is the sentence

$$\exists w_1, w_2, \dots, w_n \forall w_{n+1} \bigvee_{i=1}^n R(w_i, w_{n+1}).$$

Let  $\Delta$  be a finite subset of  $\Gamma$ , so there is an  $m \in \mathbb{N}+$  such that if  $\neg \phi_i \in \Delta$  then  $i \leq m$ . So  $\Delta \subseteq \{\theta\} \cup \{\neg \phi_i \mid 1 \leq i \leq m\}$ . Let M be the structure for L such that  $|M| = \{1, 2, 3, \ldots, m+1\}$  and

$$R^M = \{ \langle i, i \rangle \mid 1 \le i \le m+1 \}.$$

Then  $M \models \theta$  since M has a finite cover, namely  $\{1, 2, ..., m + 1\}$ . Also  $\phi_n$  fails in M for  $n \leq m$  since for any  $j_1, j_2, ..., j_n \in |M|$ ,

$$M \nvDash \bigvee_{i=1}^{n} R(i_j, k)$$

for any k from the non-empty set

$$|M| - \{j_1, j_2, \dots, j_n\} = \{1, 2, \dots, m+1\} - \{j_1, j_2, \dots, j_n\},\$$

non-empty because

$$m+1=|\{1,2,\ldots,m+1\}|>m\geq n\geq |\{j_1,j_2,\ldots,j_n\}|.$$

Hence M is a model of  $\Delta$ .

 $\therefore$  By the above Compactness Theorem  $\Gamma$  is satisfied in some structure K for L. Hence  $K \models \theta$  so by assumption K has a finite cover,  $\{a_1, a_2, \ldots, a_n\}$  say. Therefore

$$K \models \forall w_{n+1} \bigvee_{i=1}^{n} R(a_i, w_{n+1})$$

and hence

$$K \models \exists w_1, w_2, \dots, w_n \forall w_{n+1} \bigvee_{i=1}^n R(w_i, w_{n+1}),$$

i.e.  $K \models \phi_n$ . But this is a contradiction since  $K \models \Gamma$  and  $\neg \phi_n \in \Gamma$ . We conclude that no such  $\theta$  can exist, as required.