Two Hours and Thirty Minutes

UNIVERSITY OF MANCHESTER

NONSTANDARD LOGICS

24th April 2009

2.00 - 4.30

Answer all 8 questions (80 marks in all).

A list of axioms and rules of proof is appended to this examination paper

SECTION A

Answer $\underline{\mathbf{ALL}}$ 8 questions

A1. How is
$$\succ_{\vec{s}}$$
 defined for $\vec{s} = s_1, \ldots, s_m \subseteq At^L$?

State the Representation Theorem for Rational Consequence Relations. In the case when $L = \{p, q, r\}, \vec{s} = s_1, s_2, s_3$ and

$$s_1 = \{\neg p \land \neg q \land \neg r\}$$

$$s_2 = \{p \land \neg q \land \neg r, \neg p \land q \land r\}$$

$$s_3 = \{p \land q \land r\}$$

which of the following are true? [You need not justify your answers.]

(i)
$$\neg p \succ_{\vec{s}} \neg q \land \neg r$$

(ii) $p \lor q \succ_{\vec{s}} r$
(iii) $q \land \neg r \succ_{\vec{s}} p$

[12 marks]

A2. By giving a direct derivation from the GM rules show that the rule

$$\frac{\neg\phi \succ \psi \quad \neg\phi \not\sim \neg\theta}{\theta \succ \phi \lor \psi}$$

is satisfied by all rational consequence relations. [You may use the derived rule Con but in that case you should give a derivation of it from the GM rules.]

[6 marks]

A3. By using the Representation Theorem for rational consequence relations show that the rule

$$\frac{\theta \lor \phi \mathrel{\hspace{0.2em}\sim} \neg \theta}{\psi \lor \phi \mathrel{\hspace{0.2em}\sim} \neg \theta}$$

holds for all rational consequence relations, but that the rule

$$\frac{\theta \lor \phi \mathrel{\mid} \sim \neg \theta}{\psi \land \phi \mathrel{\mid} \sim \neg \theta}$$

fails for some rational consequence relation (and choice of θ, ϕ, ψ).

[12 marks]

P.T.O.

A4. How is the relation \models^{S_4} defined? State the Completeness Theorem for S_4 . Show that:

(i) $\models^{S_4} \diamondsuit \theta \lor \Box \Box \neg \theta$.

(ii) $\nvDash^{S_4} p \lor \Box \Box \neg p$.

[12 marks]

A5. What is meant by a *proof* in *B*? What does it mean to say $\Gamma \vdash^{B} \theta$? Give a formal proof to show that

$$\vdash^B \Box \theta \to \Box \Box \Diamond \theta$$

[8 marks]

A6. Let H be K together with the axiom schema

 $\Diamond \theta \mid \Diamond \Diamond \theta$

By considering a suitable family of frames, or otherwise, show that

$$\Diamond p \nvDash^H p$$

[10 marks]

A7. List the desiderata (C1)-(C4) for a function $F_{\wedge} : [0,1]^2 \to [0,1]$.

What can be said about F_{\wedge} according to the classification given by the Mostert-Shields Theorem when, in addition to (C1)-(C4), F_{\wedge} also satisfies that $0 < F_{\wedge}(x, x) < x$ for all $x \in (0, 1)$?

[9 marks]

A8. How is the relation $\models^{\underline{L}}$ defined?

Which of the following are true (for all $\theta, \phi \in SL$)? In each case you should justify your answer. (i) If $\theta \models^{L} \phi$ then $\models^{L} \theta \rightarrow \phi$. (ii) $\models^{L} ((\phi \rightarrow \theta) \rightarrow \phi) \rightarrow (\theta \rightarrow \phi)$.

[11 marks]

END OF EXAMINATION PAPER

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Axioms and Rules of Proof

Sentential, or Propositional, Calculus, SC

 $\underline{\text{Axioms:}} \quad \text{REF} \quad \theta \mid \theta \\ \underline{\text{Rules of Proof:}}$

Modal Logics

Axioms:

 $\begin{array}{lll} \mathrm{K} := & \mathrm{REF} + \Box(\theta \to \phi) \mid \Box \theta \to \Box \phi \\ \mathrm{T} := & K + \Box \theta \mid \theta \\ \mathrm{D} := & K + \Box \theta \mid \Diamond \theta \\ \mathrm{B} := & K + \theta \mid \Box \Diamond \theta \\ \mathrm{S}_4 := & T + \Box \theta \mid \Box \Box \theta \\ \mathrm{S}_5 := & T + \Diamond \theta \mid \Box \Diamond \theta \end{array}$

<u>Rules of Proof:</u> All the rules of *SC* plus NEC $\frac{|\theta|}{|\Box \theta|}$

Lukasiewicz Logic, Ł

Axioms: REF together with

L1:
$$|\theta \to (\phi \to \theta)$$

L2: $|(\theta \to \phi) \to ((\phi \to \psi) \to (\theta \to \psi))$
L3: $|(\neg \theta \to \neg \phi) \to (\phi \to \theta)$
L4: $|((\theta \to \phi) \to \phi) \to ((\phi \to \theta) \to \theta)$ i.e. $|(\theta \lor \phi) \to (\phi \lor \theta)$
L5: $|(\theta \to \phi) \lor (\phi \to \theta)$

<u>Rules of Proof:</u> Only MP

GM Rules and Axiom

REF, $\theta \succ \theta$, together with:

left logical equivalence	$\frac{\theta \equiv \phi, \theta \succ \psi}{\phi \succ \psi}$	LLE
right weakening	$\frac{\theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \phi, \hspace{0.2em} \phi \models \psi}{\theta \hspace{0.2em}\mid\hspace{0.58em} \psi}$	RWE
cautious monotonicity	$\frac{\theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \phi, \hspace{0.2em} \theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \psi}{\theta \wedge \phi \hspace{0.2em}\mid\hspace{-0.58em}\mid\hspace{0.58em} \psi}$	СМО
and on right	$\frac{\theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \phi, \hspace{0.2em} \theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \psi}{\theta \hspace{0.2em}\mid\hspace{0.58em}\sim\hspace{-0.9em} \phi \wedge \psi}$	AND
disjunction, or left	$\frac{\theta \hspace{0.2em}\sim\hspace{-0.9em}\mid\hspace{0.58em} \psi, \hspace{0.2em} \phi \hspace{0.2em}\mid\hspace{0.58em} \psi}{\theta \lor \phi \hspace{0.2em}\mid\hspace{0.58em} \psi}$	DIS
rational monotonicity	$\frac{\theta \not\succ \psi, \theta \not\succ \neg \phi}{\theta \land \phi \not\succ \psi}$	RMO